

# Counting

## Permutations and Combinations

Def Given a non-empty finite set  $B$  (known as an **alphabet**), a **string** of length  $k \in \mathbb{N}$  in  $B$  is a mapping  $\sigma: \{1, \dots, k\} \rightarrow B$ .

There exist  $n^k$  length- $k$  strings over an alphabet of size  $n$ .

This can easily be proved by induction on  $k$ .

Note that the empty string is well-defined.

A **binary string** is a string over an alphabet of size 2.

A length- $k$  binary string can be used to represent a subset of a set of size  $k$ .

More concretely, let  $[k] = \{1, 2, \dots, k\}$  and  $B = \{0, 1\}$ . Then given any  $\sigma$ , we can get a subset of  $[k]$  by  $\{i : \sigma(i) = 1\}$ .

Further note that there is a one-one mapping between subsets and strings. We often represent  $\sigma(i)$  as  $\sigma_i$ .

A **permutation** of  $B$ , an alphabet of size  $n$ , is a bijection from  $[n]$  to  $B$ .

This is just a length  $n$  string with no repeating characters.

For example,  $\begin{matrix} c & a & d & e & b \\ 1 & 2 & 3 & 4 & 5 \end{matrix}$  is a permutation of  $\{a, b, c, d, e\}$ .

If we loosen the criteria by considering any string with no repeating characters, we have a one-one function from  $[k]$  to  $B$  instead of a bijection.

Let  $P(n, k)$  represent the number of such length  $k$  strings with no repeating characters given an alphabet of size  $n$ .

If  $k > n$ , then  $P(n, k) = 0$ . Otherwise, we claim  $P(n, k) = \frac{n!}{(n-k)!}$ .  
(Here,  $n! = \begin{cases} 0, & n=0 \\ n(n-1)!, & n>0 \end{cases}$ ) This is easily proved by induction on  $n$ .  
( $P(n, k) = nP(n-1, k-1)$ )

On the other hand, how many subsets of size  $k$  does a set of size  $n$  have?

We can represent subsets as strings without repetition. ( $\{a, b, c\} = abc$ )  
Here, the same subset can be represented by multiple strings.

( $abc$  and  $bca$  represent the same set)

We know how many such strings there are! Exactly  $k!$  strings which involve those  $k$  symbols.

Therefore, the number of subsets of size  $k$  is  $\frac{P(n, k)}{k!} = \frac{n!}{k!(n-k)!}$ .

We represent this as  $C(n, k)$  or  $\binom{n}{k}$  (read "n choose k")

Note that  $C(n, k) = C(n, n-k)$ .

(choosing a subset of size  $k$  is the same as choosing a complement of size  $n-k$ )

$C(0, 0) = 1$ . This just says  $\emptyset \subseteq \emptyset$ .

$C(n, 0) + C(n, 1) + \dots + C(n, n) = 2^n$  (there are  $2^n$  subsets)

More generally,

Theo. [Binomial Theorem]

For any  $x \in \mathbb{R}$  and  $n \in \mathbb{N}_0$ ,

$$(1+x)^n = \sum_{k=0}^n C(n, k) x^k$$

↳ think of this as choosing which  $k$   $x$ 's to multiply.

This can be proved using induction. To do so, show that

$$C(n, k) = C(n-1, k-1) + C(n-1, k)$$

This can be thought of as: let  $|S| = n$  and  $a \in S$ . Count the number of subsets of size  $k$  that include  $a$  and the number that don't include  $a$ .

$$\begin{aligned} &\hookrightarrow C(n-1, k) && \hookrightarrow C(n-1, k-1) \end{aligned}$$

We now consider a series of "balls and boxes" problems.

In how many ways can you throw a set of balls into a set of boxes?

This has different answers depending on whether the balls/boxes are (in)distinguishable.

Balls Boxes	Labelled	Unlabelled
Labelled	→ Function	Multiset
Unlabelled	→ Set Partition	Integer Partition

We can also have some more variants: no box is empty  
at most one ball per box.

### Distinguishable balls/Distinguishable Boxes

This is just a function from  $A$ , the set of balls to  $B$ , the set of bins.

(each ball is thrown into a single bin)

The number of ways of throwing is just the number of functions from  $A$  to  $B$ . Such a function can be represented a string of length  $|A|$  over  $B$ .

⇒ The number of functions from  $A$  to  $B$  is  $|B|^{|A|}$ .

If every box can hold at most one ball, the function is one-one, which corresponds to a permutation of length  $|A|$  over  $B$ .

⇒ The number of one-one functions from  $A$  to  $B$  is  $P(|B|, |A|)$

If no box is empty, the function is onto.

Recall the inclusion-exclusion principle:  $|S \cup T| = |S| + |T| - |S \cap T|$ .

More generally, given finite sets  $T_1, \dots, T_n$ ,

$$\left| \bigcup_{i \in [n]} T_i \right| = \sum_{\substack{J \subseteq [n] \\ J \neq \emptyset}} (-1)^{|J|+1} \left| \bigcap_{j \in J} T_j \right|$$

Prove this by induction on  $n$ .

Without loss of generality, let  $A = [k]$  and  $B = [n]$ . We denote the number of onto functions from  $A$  to  $B$  by  $N(k, n)$

We claim

$$N(k, n) = \sum_{i=0}^n (-1)^i C(n, i) (n-i)^k$$

$\hookrightarrow n^k - C(n, 1)(n-1)^k + \dots$

The set of non-onto functions is

$$\left( \bigcup_{i \in [n]} T_i \right) \text{ where } T_i = \{f: A \rightarrow B \mid i \notin \text{Im}(f)\}$$

Use the inclusion-exclusion principle! The cardinality of this set is

$$\sum_{\substack{J \subseteq [n] \\ J \neq \emptyset}} (-1)^{|J|+1} \left| \bigcap_{j \in J} T_j \right| = \sum_{\substack{J \subseteq [n] \\ J \neq \emptyset}} (-1)^{|J|+1} (n-|J|)^k$$

$f$  is just a function  
from  $[k]$  to  $[n] \setminus J$   
 $\Rightarrow$  there are  $(n-|J|)^k$   
such functions

Given some  $|J|$ , how many such  $J$  exist? This number is equal to  $C(n, |J|)$ . Letting  $i = |J|$  on the right,

$$\Rightarrow \left| \bigcup_{i \in [n]} T_i \right| = \sum_{j=1}^n (-1)^{j+1} C(n, j) (n-j)^k$$

Therefore, the number of onto functions is  $n^k$  - this, namely,

$$N(k, n) = \sum_{i=0}^n (-1)^i C(n, i) (n-i)^k$$

### Indistinguishable balls / Distinguishable boxes

We only care about the number of balls in each box, not the identity of the balls.

A **multi-set** is like a set, but allows elements to occur multiple times.

Only multiplicity matters:  $[a, a, b] = [a, b, a]$ .

A multi-set is just a multiplicity function  $\mu: B \rightarrow \mathbb{N}_0$ .

The **size** of a multi-set is the sum of its multiplicities.

Here, the question is equivalent to finding the number of multisets of size  $k$  with the base set as  $[n]$ .

We want the number of  $(n_1, n_2, \dots, n_n)$  with  
 $n_1 + n_2 + \dots + n_n = k$  (Here  $n_n = \mu(n)$ )  
↳ multiplicity

Any such thing can be represented with  $n-1$  "bars" and  $k$  "stars".

So  $(n_1, n_2, n_3, n_4, n_5) = (1, 2, 1, 0, 1)$  is

\* | \*\* | \* || \*       $k=5, n=5$

↳ box 1      ↳ box 2      ↳ box 3 ...

There is a one-one correspondence between such "star-bar" combinations and the required result.

Any such combination can be formed by choosing some  $n-1$  positions among to  $n+k-1$  positions, filling bars there, and filling stars everywhere else.

⇒ The number of combinations is  $C(n+k-1, n-1)$

If we want each box to be non-empty, then just throw one ball into each box and solve the problem for a smaller  $n$  ( $=n-k$ )

(we get  $C(n-1, n-k-1)$ )

If at-most one ball per box, then we just want a set of size  $k$ , so there are  $C(n, k)$  possibilities.

### Distinguishable balls / Indistinguishable boxes

We partition the set  $A$  of balls into unlabelled bins.

We must just find the number of partitions of  $A$ .

↳ we defined this when studying equivalence relations

How many partitions does a set  $A$  of  $k$  elements have?

Let  $S(k, n)$  denote the number of partitions of  $[k]$  into exactly  $n$  parts.

↳ (This is the no bin empty case)  
↳ "Stirling number of the second kind"

More generally, the number of ways  $A$  can be partitioned into at most  $n$  parts is  $\sum_{m \in [n]} S(k, m)$

Further,  $B_k = \sum_{m \in [k]} S(k, m)$  is the total number of partitions of  $[k]$ .  
↳ Bell number.

Now, in  $S(k, n)$ , suppose the parts are labelled  $1, 2, \dots, n$ . There are  $N(k, n)$  such partitions. However, disregarding the labelling, each partition is counted  $n!$  times.

$$\Rightarrow S(k, n) = \frac{N(k, n)}{n!}$$

### Indistinguishable balls/Indistinguishable boxes

This problem is equivalent to writing  $k$  as a sum of  $n$  non-negative integers.

⇒ The number of  $(x_1, x_2, \dots, x_n)$  such that

$$x_1 + x_2 + \dots + x_n = k \text{ and } 0 \leq x_1 \leq x_2 \leq \dots \leq x_n$$

If no box is empty, then each  $x_i$  is positive.

↳ The number of such solutions is called the **partition number**  $p_n(k)$ .

If there is no restriction, the answer is just  $p_n(n+k)$ .

(Let  $y_i = x_i + 1$  for each  $i$ )

Let us now examine the partition number.

$$p_n(k) = \left| \{ (x_1, \dots, x_n) : x_1 + \dots + x_n = k \text{ and } 1 \leq x_1 \leq \dots \leq x_n \} \right|$$

First of all,  $p_0(0) = 1$   
 $p_0(k) = 0$  for  $k > 0$   
 $p_n(k) = 0$  if  $k < n$

We then have

$$p_n(k) = p_n(k-n) + p_{n-1}(k-1)$$

$\downarrow$  the case  $x_i > 1$                        $\downarrow$  the case  $x_i = 1$

This enables us to (recursively) define  $p_n(k)$  for every  $k$  and  $n$ .

Boxes \ Balls	Distinguishable	Indistinguishable
Distinguishable	$n^k$ $\left( \begin{array}{l} \text{if one-one, then } P(n,k) \\ \text{if onto, then } N(k,n) \end{array} \right)$	$C(n+k-1, k)$ $\left( \begin{array}{l} \text{if one-one, then } C(n,k) \\ \text{if onto, then } C(k-1, n-1) \end{array} \right)$
Indistinguishable	$\sum_{m \in [n]} S(k,m)$ $\left( \begin{array}{l} \text{if one-one, then } 0 \text{ or } 1 \\ \text{if onto, then } S(k,n) \end{array} \right)$	$p_n(n+k)$ $\left( \begin{array}{l} \text{if one-one, then } 0 \text{ or } 1 \\ \text{if onto, then } p_n(k) \end{array} \right)$

where  $N(k,n) = \sum_{i=0}^n (-1)^i C(n,i) (n-i)^k$ .

$$S(k,n) = \frac{N(k,n)}{n!}$$

$$p_n(k) = p_n(k-n) + p_{n-1}(k-1), \quad p_0(0) = 1, \quad p_0(k) = 0 \text{ for } k > 0$$

$$p_n(k) = 0 \text{ if } k < n$$